Efficient stencil computations in Accelerate

Josh Meredith, UNSW

Stencil computations

• Like map operation, but including neighbouring elements

Input:				
1	2	3	4	
5	6	7	8	
9	10	11	12	

Map:

f(1)	f(2)	f(3)	f(4)
f(5)	f(6)	f(7)	f(8)
f(9)	f(10)	f(11)	f(12)

f = function of current element

Stencil computations

• Like map operation, but including neighbouring elements

Input:				
1	2	3	4	
5	6	7	8	
9	10	11	12	

Map:

f(1)	f(2)	f(3)	f(4)
f(5)	f(6)	f(7)	f(8)
f(9)	f(10)	f(11)	f(12)

f = function of current element

Stencil:

 g()	

g = function of neighbouring elements

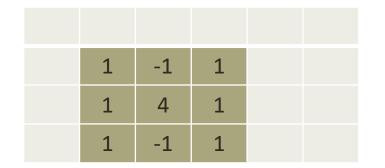
Stencil computations

- Why are they useful?
 - Image processing (e.g. photoshop filters)
 - Scientific applications (numerical simulations)

How do we implement these efficiently?

Benchmark example

High level Accelerate (Haskell) stencil:





Benchmarks (teaser)

	Time (ms)	Runtime vs C++ (x slower)
Accelerate (current)	3052	28x
C++ (hand optimised)	110	1x
This work	112	1x

Problem #1: nested loops

- Accelerate currently uses general, but slow, 1-dimensional indices
- We want an efficient implementation for 2D stencils, which are very common

Problem #1: nested loops

1-dimensional:

0	1	2	3
4	5	6	7
8	9	10	11

- i = [0, width * height)
 x = i `mod` width
- y = i `div` width

- Loop over i
- When we need x and y, calculate them with mod and div

Problem #1: nested loops

2-dimensional:

(0,0)	(0,1)	(0,2)	(0,3)
(1,0)	(1,1)	(1,2)	(1,3)
(2,0)	(2,1)	(2,2)	(2,3)

- x = [0, width)
- y = [0, height)

i = y * width + x

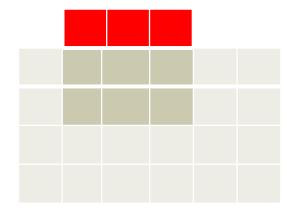
- Loop over x and y
- When we need i, calculate it based on the x, y and width

Summary problem #1

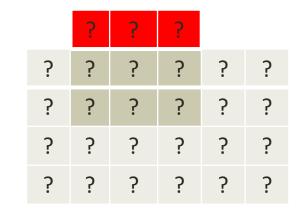
	Time (ms)	Speedup (step)	Speedup (total)
1D loop	3242	-	1

Summary problem #1

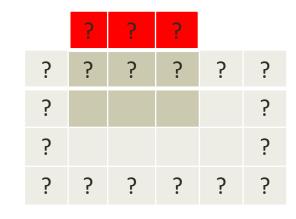
	Time (ms)	Speedup (step)	Speedup (total)
1D loop	3242	-	1
2D loop	1153	2.8	2.8



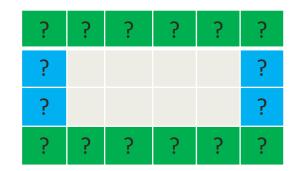
What do we do at the edges?



Naive solution: check and apply the boundary condition on all reads

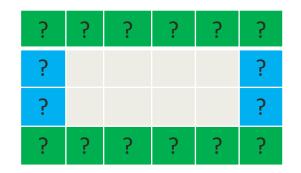


Better solution: only check the boundary condition when computing border elements



Better solution: only check the boundary condition when computing border elements

Bounds checking: corners



- Accelerate arrays are row major
- So we should include the corners with the top and bottom

Summary problem #2

	Time (ms)	Speedup (step)	Speedup (total)
1D loop	3242	-	1
2D loop	1153	2.8	2.8
Bounds checking	144	8	22.5

Problem #3: Tiling

 Reading neighbouring elements means we have a lot of duplicated reads

Solution: 1x1 tiles

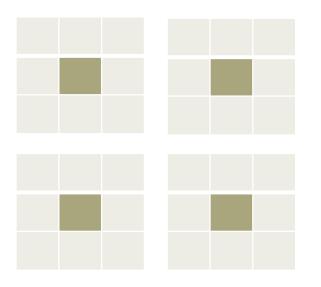
• Does LLVM make tiling unnecessary?

20

• Let's compare

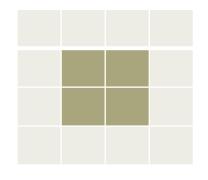
Solution: 2x2 tiles

Before tiling:



36 reads 4 writes

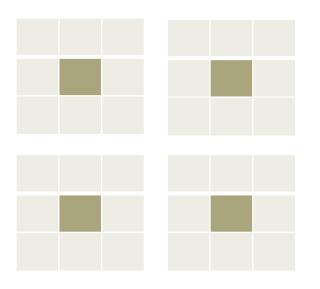
After tiling:



16 reads 4 writes

Solution: 2x2 tiles

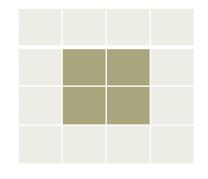
Before tiling:



36 reads 4 writes

144ms

After tiling:

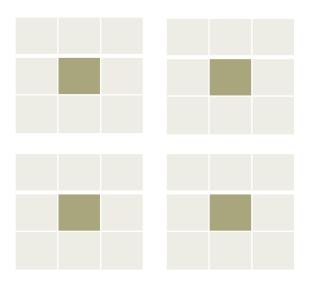


16 reads 4 writes

350ms

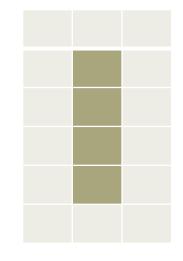
Solution: 1x4 tiles

Before tiling:



36 reads 4 writes

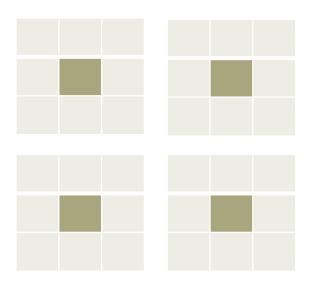
After tiling:



18 reads4 writes

Solution: 1x4 tiles

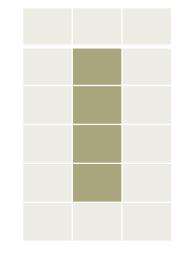
Before tiling:



36 reads 4 writes

144ms

After tiling:



18 reads 4 writes 110ms

Summary problem #3

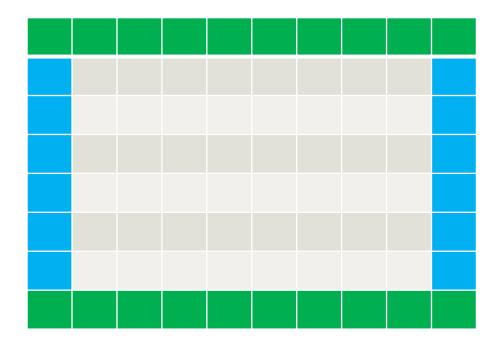
	Time (ms)	Speedup (step)	Speedup (total)
1D loop	3242	-	1
2D loop	1153	2.8	2.8
Bounds checking	144	8	22.5
1x4 tiles	110	1.3	29.5

Problem #4: Parallelisation

• How should we evaluate this in parallel?

Problem #4: Parallelisation

- Evaluate the edge regions sequentially
- Then, evaluate the main region in parallel using Accelerate's work stealer



Problem #4: Parallelisation

- No performance gains for this stencil
- Our example is very memory bottlenecked
- Many common stencils are similarly memory bottlenecked

Summary evaluation strategy

- Evaluate in several regions:
 - Top & bottom, including corners, with bounds checking
 - Left & right, excluding corners, with bounds checking
 - Middle, without bounds checking
 - Tile several elements at once on the outer axis
 - Also parallelise across the outer (y) axis
 - Vectorise across the inner (x) axis with LLVM

